Quantum Circuit Optimization by Hadamard Gate Reduction

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Abstract. Due to its fault-tolerant gates, the Clifford+T library consisting of Hadamard (denoted by H), T, and CNOT gates has attracted interest in the synthesis of quantum circuits. Since the implementation of T gates is expensive, recent research is aiming at minimizing the use of such gates. It has been shown that T-depth optimizations can be implemented efficiently for circuits consisting only of T and CNOT gates and that H gates impede the optimization significantly.

In this paper, we investigate the role of H gates in reducing the T-count and T-depth for quantum circuits. To reduce the number of H gates, we propose several algorithms targeting different steps in the synthesis of reversible functions as quantum circuits.

Experiments show the effect of H gate reductions on the costs for T-count and T-depth. Our approach yields a significant improvement of up to 88% in the final T-depth compared to the best known T-depth optimization technique.

1 Introduction

Quantum computing has shown promising results, e.g., for solving certain problems that require exponential running time in classical computers. Quantum computers exploit quantum mechanical effects and their underlying model makes use of qubits. In contrast to Boolean logic, qubits do not only represent the classical 0 and 1 states but also a superposition of both leading to a theoretically enormous speed-up in computing. The Deutsch-Jozsa algorithm [1] as well as the Shor's factorization algorithm [2] from Shor are the famous examples.

As a result, the synthesis of quantum circuits has become an active research area and many theoretical implementations for this kind of circuits have been presented [3]. To that end, since quantum operations are reversible, as a first step a reversible circuit is synthesized for the desired Boolean function after which, the resulting circuit is mapped to a functionally equivalent quantum circuit. It is also possible to build the quantum circuit for the requested Boolean function directly without going through the reversible circuit synthesis stage [4]. For the synthesis of, or mapping to, quantum circuits, several universal quantum gate libraries were introduced. One of the most used libraries is the Clifford+T library which is particularly interesting due to its fault-tolerant implementation [5].

S. Yamashita and S. Minato (Eds.): RC 2014, LNCS 8507, pp. 149–162, 2014.

After designing the quantum circuit, optimization techniques are often applied in order to produce a cheaper equivalent circuit. These optimization methods for the resulting Clifford+T circuits mainly focus on reducing the number of T gates and hence the T-depth on the resulting circuit because fault-tolerant implementations of T gates are considerably more expensive than those of the Clifford gates [6]. Thus, a couple of optimization techniques [7,8] targeting the T-depth minimization were introduced. The major obstacle facing the T-depth minimization techniques is the H gates since T gates cannot commute across such gates. For that reason, attempts for tackling this problem were either to reduce the T-depth for quantum circuits over the gate library $\{CNOT, T\}$ or to extend the same approach for the Clifford+T library circuits by optimizing the T-depth of subcircuits between the H gates boundaries [9]. To the best of our knowledge, no one has studied the effect of minimizing H gates for reducing the T-depth so far.

In this paper, we study the characteristics of H gates and show how they significantly restrain the movement of T gates and hence limit the ability to better parallelize T gates. We prove that reducing the not needed H gates leads to a more efficient minimization of the T-depth. To do so, we introduce a new methodology which aims to eliminate the H gates as a preprocessing step for improving the T-depth optimization results of quantum circuits.

The remainder of the paper is structured as follows: first the basics on reversible and quantum circuits are introduced in Sect. 2. The next section outlines the general idea. Section 4 gives a detailed description of the implementation of the presented approach, and experimental results are evaluated and interpreted in Sect. 5. The paper is concluded in Sect. 6.

2 Background

To keep the remainder of this paper self-contained, this section briefly introduces the basics on reversible circuits, quantum circuits, and the corresponding mapping from reversible to quantum circuits.

2.1 Reversible Circuits

A Boolean function $f: \mathbb{B}^n \to \mathbb{B}^n$ is said to be *reversible* if it is bijective, i.e., if each input pattern is uniquely mapped to a corresponding output pattern, and vice versa. Reversible functions can be realized by reversible circuits that consist of at least n lines. Reversible circuits are cascades of reversible gates that belong to a gate library. One gate library that is often used consists of multiple control Toffoli gates [10].

Definition 1. Given a set of variables $X = \{x_1, \ldots, x_n\}$, a multiple control Toffoli gate T(C, t) has control lines $C = \{x_{j_1}, x_{j_2}, \ldots, x_{j_l}\} \subset X$ and a target line $t \in X \setminus C$. The gate maps $t \mapsto t \oplus h(x_{j_1}, x_{j_2}, \ldots, x_{j_l})$ where h is defined as $h: (x_{j_1}, x_{j_2}, \cdots, x_{j_l}) \mapsto (x_{j_1} \wedge x_{j_2} \wedge \cdots \wedge x_{j_l})$. All remaining other lines are passed through unaltered.

(c) Reversible circuit



Fig. 1. Reversible circuitry

In [11] it has been shown that any reversible function $f: \mathbb{B}^n \to \mathbb{B}^n$ can be realized by a reversible circuit with n lines when using Toffoli gates.

(b) Multiple control Toffoli gate

Example 1. Figure 1(a) shows a Toffoli gate with two control lines. The control lines are either denoted by \bullet as depicted in Fig. 1(a) or represented by a Boolean function $h: \mathbf{x_1} = (x_{1_1}, x_{1_2}, \cdots, x_{1_l}) \mapsto (x_{1_1} \wedge x_{1_2} \wedge \cdots \wedge x_{1_l})$ as sketched in Fig. 1(b). The target line is denoted by \bigoplus . Figure 1(c) shows different Toffoli gates in a cascade forming a reversible circuit.

2.2 Quantum Circuits

(a) Toffoli gate

Instead of bits, quantum circuits manipulate qubits which can represent the classical Boolean values but also a superposition of them. A $qubit |\varphi\rangle$ is a vector $\binom{a}{b}$ where $a,b \in \mathbb{C}$ such that $|a|^2 + |b|^2 = 1$. If a = 1, then $|\varphi\rangle$ represents the classical 0, denoted $|0\rangle$, and if b = 1, then $|\varphi\rangle$ represents the classical 1, denoted $|1\rangle$.

In general, a quantum gate acting on n qubits represents a $2^n \times 2^n$ unitary matrix [12]. A matrix U is unitary if $U^{\dagger}U = UU^{\dagger} = I$ where $U^{\dagger} = (U^*)^T$ is the conjugate transpose of U. Using this gate definition many quantum mechanical effects such as superposition and entanglement can be formulated. Although Toffoli gates represent a unitary matrix, they are too general and thus not suitable for realizing quantum circuits [13]. In this paper, we make use of a gate library that is universal for quantum computation as well as its gates can be implemented in a fault-tolerant way.

Definition 2. We consider the gate library $\{H, Z, S, T, CNOT\}$ with

$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, \ Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \ S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}, \ T = \begin{pmatrix} 1 & 0 \\ 0 & e^{\frac{i\pi}{4}} \end{pmatrix}, \ \text{CNOT} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$
(1)

as the universal gate set. Note that the S^{\dagger} , T^{\dagger} , and NOT gates can be implemented with SSS, SSST, and HZH, respectively. The S and S^{\dagger} gates are square roots of the Z gate (given by the matrix in (1)). Similarly, the T and T^{\dagger} gates are given by matrices that are the fourth root of the Z gate.

A single qubit gate G(t) over the inputs $X = \{x_1, \ldots, x_n\}$ consists of a single target line $t \in X$, while a CNOT gate G(c,t) comprises, in addition, a single control line $c \in X$ with $t \neq c$.

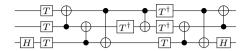


Fig. 2. Quantum circuit realizing a Toffoli gate

Example 2. Figure 2 shows a quantum circuit consisting of sixteen Clifford+T gates. This circuit represent one of the optimal realization of a Toffoli gate as depicted in [7, Fig. 13]. The circuit has a T-count of 7, a T-depth of 3, and an H-count of 2.

3 General Idea

In this work, we propose an optimization approach that aims for reducing the T-depth in a given quantum circuit as a main goal. In this section we motivate the impact of H gates on minimizing the T-depth in quantum circuits, afterwards we outline the proposed approach.

Mapping reversible circuits to quantum circuits can be done with different quantum library gates such as the NCV [12] or NCV- $|v_1\rangle$ [14] libraries. But recently the Clifford+T gate library has attracted most attention since it is composed of fault-tolerant logical gates [5]. Because it has been demonstrated that the fault-tolerant implementation of the T gates is surpassing the cost of the Clifford gates [6], many works have addressed the optimization of quantum circuits by minimizing the T-count [8] and the T-depth [9,15].

The algorithm presented in [8] describes a method that performs an exhaustive search for a circuit that implements an n-qubit unitary matrix U using the minimal number of T gates. The work introduced in [7] addressed the optimization of T-depth for small circuits composed of four qubits at maximum. This is done by applying an exhaustive search algorithm to find the optimal T-depth realization. Another approach proposed a polynomial run-time algorithm for reducing the T-depth and the T-count of quantum circuits over the gate library $\{CNOT, T\}$ [9]. The algorithm deletes redundant T gates by computing the total phase and parallelizing the T gates through Matroid partitioning. The idea is based on decomposing a given function into minimal number of linear Boolean functions and then resynthesize each one with an optimal T-depth realization. This algorithm is extended to circuits built with the Clifford+T library. In this case, the same approach is applied for the subcircuits between the H gates, afterwards an optimization process is applied which detects the identical gates and deletes them.

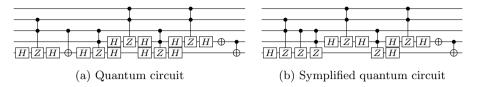


Fig. 3. Equivalent quantum circuit for Fig. 1(c)

So far, as it is explained above, all the optimization techniques proposed algorithms for improving the T-depth for quantum circuits in the absence of, or locally between, H gates. But no work has introduced an approach for optimizing quantum circuits including H gates since they present the bottleneck for finding the optimal T-depth, i.e., they cannot interact with neighbouring gates and thus block the movement of any other gates across them. This restricts possible rearrangements of T gates and hence reduces the ability to perform T gates in parallel or apply possible reduction rules to a target circuit. Also when considering the algorithm [9] explained above and taking into account that the H gates are reduced before, this will allow to have larger subcircuits to resynthesize comparing to the first subcircuits: thus we get a bigger chance to get more parallel T gates and hence lower T-depth. Following the previous observations, therefore we present an approach that minimizes the H-count of quantum circuits, enabling better optimization results for the T-depth and the T-count.

Example 3. Figure 3(a) depicts the equivalent quantum circuit for the reversible circuit drawn in Fig. 1(c) according to [9]. Following their algorithm, the circuit is partitioned into a set of subcircuits located between the H gates. As it is shown in Fig. 3(a) we have 7 subcircuits. Next step, each subcircuit is resynthesized with an optimal T-depth quantum circuit. The application of the algorithm to the quantum circuit gives a circuit with a T-depth equal to 11 and a T-count equal to 27. However, one can reduce the H-count which yields an equivalent circuit depicted in Fig. 3(b). Applying the same algorithm to this circuit, which has 5 subcircuits, results in a circuit with a T-depth of 9 and a T-count of 23.

4 Optimization Approaches

Motivated by the idea outlined in the previous section, we propose a design flow (depicted in Fig. 4) for the synthesis of a reversible function realized using gates from the Clifford+T library. First, the desired function is realized as a reversible

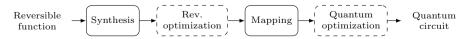


Fig. 4. Design flow for quantum circuits

circuit with Toffoli gates by applying existing synthesis methods such as [16-18]. To achieve better H gate reductions, we have taken the benefits of existing work aiming to optimize reversible circuits; for instance [19,20], and applied the template matching technique introduced in [20] to the reversible circuit. Afterwards, we have incorporated an alternative mapping technique that yields circuits which are particularly suitable for H gate reductions. Finally, the obtained quantum circuit is optimized by applying an algorithm that aims at the T-depth optimization based on H gate minimizations. To resume, although we have employed the existing optimization algorithms at the reversible optimization level, our main contributions are on the mapping and quantum optimization steps.

4.1 Optimizations at the Reversible Level

In order to enhance the obtained circuit from a synthesis approach, a post synthesis process, also called reversible optimization stage, is applied. There are many existing methods targeting the optimization of reversible circuits for either reducing the number of lines [21], number of gates [20], depth [22], or quantum cost [23]. We are interested in techniques that lead to lower quantum cost. Smaller circuits are likely to have less H-count and therefore would have better T-depth but this is not always guaranteed.

Among the interesting work that focus on reducing the quantum cost for a given reversible circuit are template matching algorithms as described in [20], the window optimization introduced in [23], and finally the algorithm outlined in [24], that is similar to the template matching algorithm but includes better gate movement properties in the whole circuit. For this work, one can apply all of these approaches along with any other method leading to optimized quantum cost. For our experiments, we have included only the template matching approach [20].

4.2 Optimizations in the Mapping

After realizing and optimizing the reversible circuit for a given reversible function, each reversible gate is mapped to its equivalent quantum circuit as described in [13]. This mapping strategy is optimized with respect to quantum cost [25]. Afterwards, a second mapping technique that leads to an even lower quantum cost was described in [26]. Another functional mapping algorithm was presented in [27]: the described method searches for gates that have the same controls but different targets and decomposes them with a special decomposition.

According to Lemma 7.3 in [13], a reversible Toffoli gate with c controls (where $c \geq 3$) can be mapped to a network consisting of two identical gates with m controls and two other identical gates with c-m+1 controls, where $m \in \{2, \cdots, c-2\}$ and each of them are placed alternately. One has a lot of freedom on how to choose the controls and the order for each gate. As an example, Fig. 5(b) presents a possible mapping for the circuit depicted in Fig. 5(a) where the partitioning of controls is done with respect to their order in the original gates. However an alternative application of Lemma 7.3 [13] results in a circuit

with two identical adjacent gates which can be removed as shown in Fig. 5(c). By removing these gates, at least two H gates are eliminated.

Hence our approach aims to apply a special mapping technique that is particularly suitable for circuits in which H gates cancel. This technique, instead of mapping reversible gates one by one and each on his own side, gathers gates as shown in the schemas in Fig. 6 and finds a suitable partitioning of the controls that leads to reversible gates that cancel and thus reduces the H-count.

This mapping technique can be applied when a pair of reversible gates have one of the structures explained as below:

- Gates having a structure similar to the Peres gates as sketched in Fig. 6(a), i.e., a control line of the first is a target line of the other, besides they share one or more controls and the first gate has its target in a non shared line.
- Gates having their targets in the same line and sharing one or more control lines as depicted in Fig. 6(b).
- Gates having their target in non shared lines as described in Fig. 6(c). Also they have one or more control lines in common.
- Gates having a structure similar to the swap gates but also they share one or more control lines as outlined in Fig. 6(d).

Example 4. Consider the case of the pair of gates depicted in Fig. 5(a). Using the new mapping scheme, we obtain the circuit drawn in Fig. 5(c) that contains 2 identical reversible gates with 3 controls each. The removal of these gates will lead to a reduction of 16 H gates compared to the classical mapping algorithm.

4.3 Optimizations at the Quantum Level

There are many optimization schemes that aim for quantum cost reduction for circuits based on the NCV library. In particular, the application of quantum template matching [28] or the merging and deletion rules together with functional moving rules as explained in [24] are beneficial for decreasing the quantum cost. However, post mapping optimization techniques designed for quantum circuits based on Clifford+T gates are limited to the reduction of identical gates as described in [9] and the identities shown in [29].

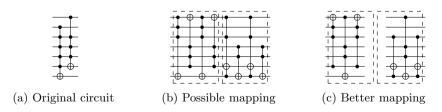


Fig. 5. Reversible circuit mapping

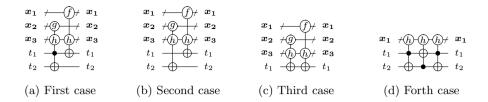


Fig. 6. Functional mapping

The optimization approach that we introduce is based on a greedy algorithm that traverses repeatedly the circuit and looks for any possible cascade replacement with a cheaper equivalent cascade or any identity deletion. These two operations are known as merging and deleting rules. This scheme can additionally be improved by applying the moving rules for quantum circuits. In fact, in the Clifford+T library, additionally to the moving rules defined in [28], a CNOT gate $G(c, t_1)$ and a root of Z gate $(Z, S, T, S^{\dagger}, T^{\dagger})$ $G(t_2)$ can be interchanged if $t_1 \neq t_2$ as it is sketched with all other possible moving rules in Fig. 7. Furthermore, the following moving and deletion rules can be exploited for the Clifford+T circuits:

Hadamard Gates Reduction. The circuit is mapped in order to locate identical H gates or one of the cascades sketched in Fig. 8. Identical gates that could be moved together are deleted from the circuit and other identified templates are replaced by its equivalent cascade that do not contain any H gate.

Merging and Deleting Gates Reduction. Taking the benefits of the moving properties for Clifford+T gates depicted on Fig. 7, additional reductions are possible for the remaining gates of the library. The algorithm searches for the templates shown in Fig. 10 and replaces these by their cheaper realization.

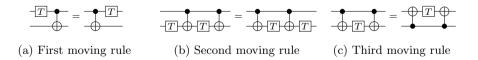


Fig. 7. Moving rules for the Clifford+T gates

Fig. 8. Reduction rules for the H gates

Fig. 9. Reduction rules for the remaining gates

5 Experimental Results

In this work, we proposed considerations of H gate minimizations to optimize quantum circuits build using the Clifford+T library. We have observed that eliminating H gates often leads to quantum circuits with a much smaller T-depth. Motivated by this, we introduced an improved design flow that aims at having lower H-count when generating the corresponding quantum gate cascades. The proposed idea described above has been implemented in the open source toolkit RevKit [30]. The experimental evaluation has been carried out on an Intel Core is Processor with 4 GB of main memory using the benchmarks taken from [31,32] database.

To determine the best synthesis approach with respect to *T-depth*, we have generated for each benchmark its corresponding circuits utilizing the following synthesis approaches: the transformation based synthesis approach (TBS [16]), the Reed-Muller synthesis approach (RMS [33]), the Young subgroups based synthesis approach (YSG [34]), and the ESOP based synthesis approach [17]. Due to space constraints, we have not detailed results for the ESOP based synthesis approach.

The experimental results are shown graphically in the plots in Fig. 10. The values of x-axis and the y-axis (logarithmic scale) denote the benchmark and the T-depth, respectively. Each plot contains three different scenarios: the T-depth of the original quantum circuits, the T-depth of the optimized circuits based on [9], and the T-depth of the optimized circuits based on our technique. One can clearly see that the T-depth and the H-count related to each other. Besides, most of the cases, the Reed-Muller synthesis approach [33] outperforms the other synthesis techniques in terms of producing lower T-depth circuits. The same observations are found for the H-count as it is shown in Fig. 10. In the rest of the paper we consider only the results of the Reed-Muller synthesis approach.

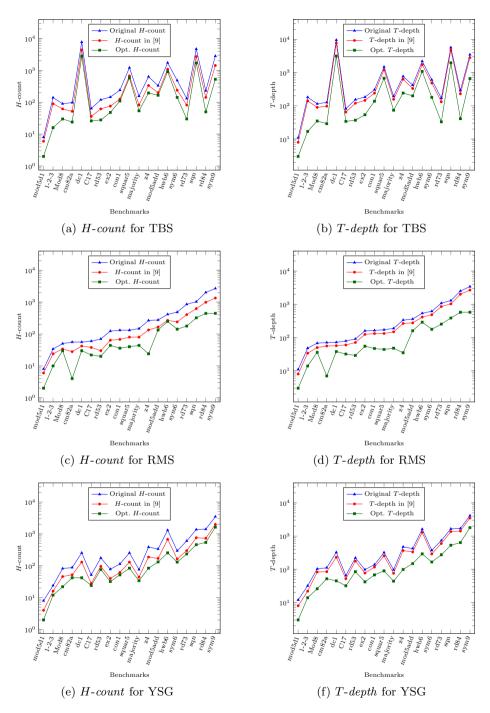


Fig. 10. H-count and T-depth for original benchmark, optimized T-depth benchmarks, optimized H-count and T-depth benchmarks

Table 1. Experimental evaluation for Reed-Muller synthesized circuits where L, QC, HC, TC, TD, Time, ΔH , and TD Imp. refer to the number of lines, the quantum costs, the H-count, the T-count, the T-depth, the needed run-time, the H gate reductions, and the T-depth improvement with regret to the costinities of invariety based on [0] respectively.

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	1035					6.870	-20	30%
$50 ext{ } 175 ext{ } 68 ext{ } 413 ext{ } 34$	131					0.032	4-	28%

Table 1 summarizes the obtained result. All benchmarks are listed in the first column. Then, the number of lines (L), the quantum costs (QC), the H-count (HC), the T-count (TC), and the T-depth (TD) of the respective circuit realizations as well as the needed run-times (Time) are provided.

The H gate reductions and the relative T-depth improvement of the circuits obtained by the proposed technique with respect to the optimized circuits based on the approach presented in [9] are provided in the columns denoted by ΔH and TD Imp., respectively.

In total three different aspects are studied: (1) the results of circuits generated from the Reed-Muller synthesis approach, (2) the results of optimized circuits using the algorithm introduced in [9], and (3) the results of the optimized circuits using the technique reviewed in Sect. 4 in addition to the approach. in [9].

Applying the T-depth optimization approach described in [9] reduces the T-depth significantly. However, it is clearly observed that these results can be improved when applying the approach based on H gate reductions. Our proposed approach leads to additional T-depth reductions of 10% in average. The results confirm the impact of eliminating H gates on the T-depth.

As can be seen, in particular our scheme leads to significant H-count reductions. Over all circuits, reductions up to 900 H gates can be obtained, therefore, this enables further improvements of the overall T-depth as it is shown in Table 1. The T-depth is reduced by 60% on average and in the best case (cm82) by 88%.

H-count and T-depth are related to each other and more the H gates are reduced, more the T-depth is lower. This explains the variation of the T-depth improvement for each benchmark. For example, when a circuit contains many Toffoli gates that have their targets in the same line, then after quantum mapping the majority of H gates will be at the same line and many will cancel. Therefore the T-depth is reduced significantly (Z4, Sym9). Whereas, when Toffoli gates have their targets in different lines then the H-count cannot be much improved, hence the T-depth is not much decreased (Mod8).

6 Conclusion

In this paper we introduced a scheme for optimizing the T-depth of quantum circuits based on H gate reductions. To that end, we incorporated a 3-level strategy targeting the optimization of circuits at the reversible, mapping and quantum level to achieve better H-count reductions and hence possible further T-depth improvements. Experimental results have shown significant T-depth reductions which reach over 80% for quantum circuits.

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